## Randomised Algorithms Supervision 4 Spectral graph theory

## 1 Linear Algebra

Exercise 1 [Properties of dagger] Prove the following properties of the conjugate transpose, defined for any $x \in \mathbb{C}^{n}$ as

$$
x^{\dagger}=\left(x_{1}^{*}, \ldots, x_{n}^{*}\right)
$$

(a) For any two vectors $x, y \in \mathbb{C}^{n}$, it holds that $x^{\dagger} y=\left(y^{\dagger} x\right)^{\dagger}$.
(b) For any vector $x \in \mathbb{C}^{n}$, it holds that $x^{\dagger} x \geq 0$.

Extended Note 1 A matrix $A \in \mathbb{C}^{n \times n}$ is Hermitian if $A^{\dagger}=A$. The notation $A^{\dagger}$ means the complex conjugate of $A$, i.e., $A^{\dagger}=\left(A^{T}\right)^{*}$. If you prefer you can attempt the following exercises, by assuming that $A$ is a matrix with real entries and then $A^{\dagger}=A^{T}=A$, meaning that the matrix is symmetric.

Exercise 2 [Real eigenvalues] Consider a Hermitian matrix $A \in \mathbb{R}^{n \times n}$, then all its eigenvalues are real.

Exercise 3 [Orthogonal eigenvectors] Consider a Hermitian matrix $A \in \mathbb{C}^{n \times n}$ and let $x, y \in \mathbb{C}^{n}$ be two eigenvectors corresponding to different eigenvalues $\lambda \neq \lambda^{\prime}$. Then, $x \perp y$.

Exercise 4 [Spectral theorem] Let $A \in \mathbb{R}^{n \times n}$ be a symmetric matrix.
(a) Let $k \leq n-1$ and let $x_{1}, \ldots, x_{k}$ be orthogonal eigenvectors of $A$. Then, there exists an eigenvector $x_{k+1}$ that is orthogonal to $x_{1}, \ldots, x_{k}$.
(b) Prove the spectral theorem,
(c) Argue that $A$ can be written as $X D X^{-1}$ for some matrix $X$ and a diagonal matrix $D$.

Exercise 5 [Inverse matrix] Let $A \in \mathbb{R}^{n \times n}$ be a symmetric matrix with (orthonormal) eigenvectors $x_{1}, \ldots, x_{n}$ corresponding to eigenvalues $\lambda_{1}, \ldots, \lambda_{n}$. Prove that:
(a) Let $x \in \mathbb{R}^{n}$ be arbitrary. By writing $x=\left(x^{T} x_{1}\right) x_{1}+\ldots+\left(x^{T} x_{n}\right) x_{n}$, show that

$$
x_{1} x_{1}^{T}+\ldots x_{n} x_{n}^{T}=I
$$

(b) By writing $A x=A I x$, show that

$$
A=\lambda_{1} x_{1} x_{1}^{T}+\ldots \lambda_{n} x_{n} x_{n}^{T}
$$

(c) Show that if $\lambda_{1} \neq 0, \ldots, \lambda_{n} \neq 0$, then

$$
A^{-1}=\frac{1}{\lambda_{1}} x_{1} x_{1}^{T}+\ldots+\frac{1}{\lambda_{n}} x_{n} x_{n}^{T}
$$

Exercise 6 [Power of a matrix] Consider a real symmetric matrix $A \in \mathbb{R}^{n \times n}$ with eigenvectors $x_{1}, \ldots, x_{n}$ and eigenvalues $\lambda_{1}, \ldots, \lambda_{n}$. Then, for any $k \in \mathbb{N}_{\geq 1}$, the eigenvectors of $A^{k}$ are $x_{1}, \ldots, x_{n}$ and the eigenvalues $\lambda_{1}^{k}, \ldots, \lambda_{n}^{k}$.

Exercise 7 [Trace of a matrix] The trace of a matrix is defined as

$$
\operatorname{tr}(A)=\sum_{i=1}^{n} A_{i i}
$$

(a) Show that for any two matrices $X$ and $Y$, we have that $\operatorname{tr}(X Y)=\operatorname{tr}(Y X)$.
(b) Show that for any real symmetric matrix $A, \operatorname{tr}(A)=\sum_{i=1}^{n} \lambda_{i}$.

Exercise 8 [Determinant of symmetric matrix] Consider a real symmetric matrix $A$. Show that

$$
\operatorname{det}(A)=\prod_{i=1}^{n} \lambda_{i}
$$

Hint: Use the property that $\operatorname{det}(A B)=\operatorname{det}(A) \cdot \operatorname{det}(B)$.

## 2 Graph matrices

### 2.1 Adjacency matrix

Exercise 9 [Basic properties] Consider the adjacency matrix of an undirected graph.
(a) Show that $\operatorname{deg}\left(v_{i}\right)=\sum_{j=1}^{n} A_{i j}$.
(b) Show that the adjacency matrix of a $d$-regular graph has eigenvalue $d$.

Exercise 10 [Counting paths] Consider the adjacency matrix $A$ of a graph $G$.
(a) Show that if $\left(A^{k}\right)_{i j}>0$ for some integer $k>0$ then there is a path of length $k$ connecting $i$ and $j$.
(b) Show that $\left(A^{k}\right)_{i j}$ also gives the number of paths connecting $i$ and $j$ with $k$ hops.
(c) Interpret $\operatorname{tr}\left(A^{k}\right)>0$.

Exercise 11 [Bipartite graphs] Show that for any bipartite graph $G$ with adjacency matrix $A$, if $\lambda>0$ is an eigenvalue then $-\lambda$ is also an eigenvalue.

## Exercise 12 [Converse for bipartite graphs]

(a) Argue that a graph with no odd length cycle is bipartite.
(b) Show that the statement "a graph has no odd length cycle" is equivalent to "for every odd $k$, $\operatorname{tr}\left(A^{k}\right)=0 "$.
(c) Deduce that if for every eigenvalue $\lambda$ of $A$, there is another eigenvalue $-\lambda$, then the graph is bipartite.

Exercise 13 On Lecture 11/slide 10 (Example 1) we determined the spectrum of the adjacency matrix $A$ for the complete graph (a.k.a. clique) of size 3. Here we would like to generalise this to any complete graph of size $n>3$. Prove that the spectrum consists of eigenvalues $n-1$ (with multiplicity 1 ) and -1 (with multiplicity $n-1$ ).

Exercise 14 [Perron-Frobenius] Let $G$ by a connected graph with adjacency matrix $A$ with eigenvectors $x_{1}, \ldots, x_{n}$ and eigenvalues $\lambda_{1}, \ldots, \lambda_{n}$, then show that
(a) $\lambda_{1} \geq-\lambda_{n}$,
(b) $\lambda_{1}>\lambda_{2}$,
(c) There exists an eigenvector $x_{1}$ which has all its entries $>0$.

### 2.2 Laplacian matrix

Exercise 15 [Factorisation] Consider the unormalised Laplacian matrix $L=D-A$ and the incident $\operatorname{matrix} M \in \mathbb{R}^{n \times m}$ defined as $M_{u e}=\mathbf{1}_{u \in e}$ (i.e., indicates which edges contain which vertices). Show that

$$
L=M^{T} M
$$

Exercise 16 Consider an undirected, $d$-regular graph $G$ and the matrices $A_{G}$ and $L_{G}$.
(a) Show that the two matrices have the same eigenvectors.
(b) Describe the correspondence between their eigenvalues.

Exercise 17 Consider a $d$-regular graph and its Laplacian matrix $L$.
(a) Using the quadratic form, show that for any vector $x \in \mathbb{R}^{n}($ with $x \neq 0)$,

$$
\frac{x^{T} L x}{x^{T} x} \leq 2
$$

(b) Deduce that $\lambda_{n} \leq 2$.

Prove that for any $d$-regular graph, the largest eigenvalue of the Laplacian $L$ satisfies $\lambda_{n} \leq 2$.

Exercise 18 Show that if $G$ is an undirected, $d$-regular, connected and bipartite graph, then the largest eigenvalue $\lambda_{n}$ of the Laplacian matrix satisfies $\lambda_{n}=2$ (this proves one direction of the fourth statement in the Lemma from Lecture 11/slide 14).

Exercise 19 Redo Exercise 19 without assuming that the graph is not $d$-regular.

Exercise 20 Consider the transition matrix of a lazy random walks $\tilde{P}=(P+I) / 2$ on a $d$-regular graph (here $I$ is the $n \times n$ identity matrix and $P$ is the transition matrix of a simple random walk).
(a) Using Exercise 17 and that the eigenvalues of $L$ are in [0,2], argue that the eigenvalues of $A$ are in $[-d, d]$.
(b) Prove that all eigenvalues of $\tilde{P}$ are non-negative.

## 3 Conductance

## Exercise 21 [Conductance of graphs]

(a) Compute the conductance of the complete graph $K_{n}$.
(b) Compute the conductance of the cycle $C_{n}$.
(c) Compute the conductance of a path $P_{n}$.
(d) Compute the conductance of a 2D grid.
(e) $(+)$ Compute the conductance of a 3D grid.

## Exercise 22

(a) Prove that for every $n>2$ there is an unweighted, undirected $n$-vertex graph with conductance 1 .
(b) (+) Can you characterise all graphs with that property?

Exercise 23 Prove that for any $d$-regular graph with $n \rightarrow \infty$ being large, the conductance satisfies $\Phi(G) \leq \frac{1}{2}+o(1)$.
Hint: Use the probabilistic method to construct a set $S$ with the required conductance. First obtain bounds for $||S|-n / 2|$ and then for $\left|E\left(S, S^{c}\right)-|E| / 2\right|$.

Exercise 24 Consider an undirected, and $d$-regular graph $G=(V, E)$ with conductance $\Phi>0$. In this exercise, you will show that the diameter of the graph is at most $\mathcal{O}\left(\frac{\log n}{\Phi}\right)$.
(a) Consider an arbitrary vertex $u \in V$ and let $S_{0}:=\{u\}$ and $S_{i}:=B_{\leq i}(u)$, the set of nodes at a distance of at most $i$ from $u$. Show that for any $S_{i}$ with $\left|S_{i}\right| \leq n / 2$, it holds that

$$
\left|E\left(S_{i}, S_{i}^{c}\right)\right| \geq \Phi \cdot\left|S_{i}\right| \cdot d
$$

and that

$$
\left|S_{i+1}\right| \geq\left|S_{i}\right| \cdot(1+\Phi)
$$

(b) Using that $\log (1+z) \geq(1 / 2) \cdot z$ (for any $z \in[0,1]$ ), deduce that $\left|S_{i}\right|>n / 2$ for $i>2 \cdot \frac{\log n}{\Phi}$.
(c) Deduce that there is no pair of vertices at a distance $>\frac{4 \log n}{\Phi}$.

